

Prolog Parsing Assistant for Categorical Grammar*

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*This presentation was prepared with L^AT_EX, in its Linux implementation — TeT_EX— and with an IDE program — Kile — both installed in computers running on a Brazilian distribution of Linux — Kurumin — everything according to the free software guidelines.

1 Introduction

- Quasi-arithmetical notation for linguistic description
- AB-model (Ajdukiewicz & Bar-Hillel) is weakly equivalent to context-free grammar
- Classical Categorical Grammars are mild context-sensitive
- Lambek calculus: generalization by natural deduction
- Structural completeness: “for any syntactically coherent string, any substring can be derived as a constituent and any item in that substring as its head” [2, p. 5]

2 Categorical Grammars

- function and argument relation rather than constituent structure
- syntax and semantics homomorphism
- monotonicity (there are no destructive rules such as movement or deletion)
- radical lexicalism

3 Free Categorical Grammar [1]

3.1 Application (R1)

- $X/Y : \phi \quad Y : \alpha \rightarrow X : (\phi \alpha)$
- $Y : \alpha \quad Y \backslash X : \phi \rightarrow X : (\phi \alpha)$

Pedro	ama	Maria
$\frac{}{N : p} \quad Lex$	$\frac{}{(N \backslash S)/N : A} \quad Lex$	$\frac{}{N : m} \quad Lex$
	$\frac{}{N \backslash S : (A \ m)} \quad R1$	
	$\frac{}{S : ((A \ m) \ p)} \quad R1$	

3.2 Associativity (R3)

- $(X \setminus Y) / Z : \phi \rightarrow X \setminus (Y / Z) : \lambda x_m. \lambda x_n. ((\phi x_n) x_m)$
- $X \setminus (Y / Z) : \phi \rightarrow (X \setminus Y) / Z : \lambda x_m. \lambda x_n. ((\phi x_n) x_m)$

$[x_m \text{ and } x_n \text{ fresh}]$

3.3 β -reduction

- $(\lambda x_n. \alpha \beta) \rightarrow \alpha^{x_n \mapsto \beta}$

Pedro		ama	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S) / N : A}$	<i>Lex</i>
		$\frac{}{N \backslash (S / N) : \lambda x_0. \lambda x_1. ((A \ x_1) \ x_0)}$	<i>R3</i>
		$\frac{}{S / N : (\lambda x_0. \lambda x_1. ((A \ x_1) \ x_0) \ p)}$	<i>R1</i>
		$=_{\beta-red.} \lambda x_1. ((A \ x_1) \ p)$	

3.4 Composition (R2)

- $X/Y : \phi \quad Y/Z : \gamma \rightarrow X/Z : \lambda x_n.(\phi (\gamma x_n))$
- $Z \setminus Y : \gamma \quad Y \setminus X : \phi \rightarrow Z \setminus X : \lambda x_n.(\phi (\gamma x_n))$

$[x_n \text{ fresh}]$

3.5 Raising (R4)

- $X : \alpha \rightarrow Y/(X \setminus Y) : \lambda x_n.(x_n \alpha)$
- $X : \alpha \rightarrow (Y/X) \setminus Y : \lambda x_n.(x_n \alpha)$

$[x_n \text{ fresh}]$

Pedro		ama
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \setminus S)/N : A}$
$\frac{}{S/(N \setminus S) : \lambda x_0.(x_0 p)}$	<i>R4</i>	
$\frac{}{S/N : \lambda x_1.(\lambda x_0.(x_0 p) (A x_1))}$	<i>R2</i>	
		$=_{\beta-red.} \lambda x_1.((A x_1) p)$

4 Parsing Assistant (sample session)

```
$ swipl
Welcome to SWI-Prolog...
?- [agcl].
```

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*-----*
| Assistente de Análise Gramatical |
|           para                   |
| Gramática Categorical Livre      |
*-----*
|      Luiz Arthur Pagani         |
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*-----*
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Expressão: [pedro,ama,maria].
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1.

$n : p$, por lx
pedro

2.

$(n \backslash s) / n : A$, por lx
ama

3.

$n : m$, por lx
maria

Análise: $r4(3, \backslash)$.

Pedro		ama		Maria
$\frac{\quad}{N : p}$	Lex	$\frac{\quad}{(N \backslash S) / N : A}$	Lex	$\frac{\quad}{N : m}$
				Lex

1.

$n : p, \text{ por } lx$
 pedro

2.

$(n \backslash s) / n : A, \text{ por } lx$
 ama

3.

$(y_0 / n) \backslash y_0 : \lambda x_0. (x_0 \ m), \text{ por } r4$
 $n : m, \text{ por } lx$
 maria

Análise: r3 2.

Pedro		ama		Maria	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S) / N : A}$	<i>Lex</i>	$\frac{}{N : m}$	<i>Lex</i>
				$\frac{}{(y_0 \backslash N) / y_0 : \lambda x_0. (x_0 \ m)}$	<i>R4</i>

1.

$n : p, \text{ por } lx$
 pedro

2.

$n \backslash (s/n) : \lambda x_1. \lambda x_2. ((A \ x_2) \ x_1), \text{ por } r3$
 $(n \backslash s)/n : A, \text{ por } lx$
 ama

3.

$(y_0/n) \backslash y_0 : \lambda x_0. (x_0 \ m), \text{ por } r4$
 $n : m, \text{ por } lx$
 maria

Análise: r2 3.

Pedro		ama		Maria	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S)/N : A}$	<i>Lex</i>	$\frac{}{N : m}$	<i>Lex</i>
		$\frac{}{N \backslash (S/N) : \lambda x_1. \lambda x_2. ((A \ x_2) \ x_1)}$	<i>R3</i>	$\frac{}{(y_0 \backslash N)/y_0 : \lambda x_0. (x_0 \ m)}$	<i>R4</i>

1.

$n : p$, por lx
pedro

2.

$n \backslash s : \lambda x_3.(\lambda x_0.(x_0 \ m) (\lambda x_1.\lambda x_2.((A \ x_2) \ x_1) \ x_3))$, por r2

$n \backslash (s/n) : \lambda x_1.\lambda x_2.((A \ x_2) \ x_1)$, por r3

$(n \backslash s)/n : A$, por lx

ama

$(s/n) \backslash s : \lambda x_0.(x_0 \ m)$, por r4

$n : m$, por lx

maria

Análise: red 2.

Pedro		ama		Maria	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S)/N : A}$	<i>Lex</i>	$\frac{}{N : m}$	<i>Lex</i>
		$\frac{}{N \backslash (S/N) : \lambda x_1.\lambda x_2.((A \ x_2) \ x_1)}$	<i>R3</i>	$\frac{}{(S \backslash N)/S : \lambda x_0.(x_0 \ m)}$	<i>R4</i>
		$\frac{}{N \backslash S : \lambda x_3.(\lambda x_0.(x_0 \ m) (\lambda x_1.\lambda x_2.((A \ x_2) \ x_1) \ x_3))}$			<i>R2</i>

(β -reduction $\times 3$)

1.

$n : p$, por lx
pedro

2.

$n \backslash s : \text{lambda-}x_3.((A \ m) \ x_3)$, por redução-beta
 $n \backslash s : \text{lambda-}x_3.(\text{lambda-}x_2.((A \ x_2) \ x_3) \ m)$, por redução-beta
 $n \backslash s : \text{lambda-}x_3.(\text{lambda-}x_1.\text{lambda-}x_2.((A \ x_2) \ x_1) \ x_3) \ m)$, por redução-beta
 $n \backslash s : \text{lambda-}x_3.(\text{lambda-}x_0.(x_0 \ m) \ (\text{lambda-}x_1.\text{lambda-}x_2.((A \ x_2) \ x_1) \ x_3))$, por $r2$
 $n \backslash (s/n) : \text{lambda-}x_1.\text{lambda-}x_2.((A \ x_2) \ x_1)$, por $r3$
 $(n \backslash s)/n : A$, por lx
 ama
 $(s/n) \backslash s : \text{lambda-}x_0.(x_0 \ m)$, por $r4$
 $n : m$, por lx
 maria

Análise: $r1 \ 2$.

Pedro		ama		Maria	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S)/N : A}$	<i>Lex</i>	$\frac{}{N : m}$	<i>Lex</i>
		$\frac{}{N \backslash (S/N) : \lambda x_1.\lambda x_2.((A \ x_2) \ x_1)}$	<i>R3</i>	$\frac{}{(S \backslash N)/S : \lambda x_0.(x_0 \ m)}$	<i>R4</i>
		$\frac{}{N \backslash S : \lambda x_3.(\lambda x_0.(x_0 \ m) \ (\lambda x_1.\lambda x_2.((A \ x_2) \ x_1) \ x_3))}$			<i>R2</i>
		$=_{\beta-red.} \lambda x_3.((\lambda x_1.\lambda x_2.((A \ x_2) \ x_1) \ x_3) \ m)$			
		$=_{\beta-red.} \lambda x_3.(\lambda x_2.((A \ x_2) \ x_3) \ m)$			
		$=_{\beta-red.} \lambda x_3.((A \ m) \ x_3)$			

```

1.
s : (lambda-x_3.((A m) x_3) p), por r1
  n : p, por lx
    pedro
n\s : lambda-x_3.((A m) x_3), por redução-beta
  n\s : lambda-x_3.(lambda-x_2.((A x_2) x_3) m), por redução-beta
    n\s : lambda-x_3.((lambda-x_1.lambda-x_2.((A x_2) x_1) x_3) m), por redução-beta
      n\s : lambda-x_3.(lambda-x_0.(x_0 m) (lambda-x_1.lambda-x_2.((A x_2) x_1) x_3)), por r2
        n\((s/n) : lambda-x_1.lambda-x_2.((A x_2) x_1), por r3
          (n\s)/n : A, por lx
            ama
        (s/n)\s : lambda-x_0.(x_0 m), por r4
          n : m, por lx
            maria

```

Análise: red 1.

Pedro		ama		Maria	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S) / N : A}$	<i>Lex</i>	$\frac{}{N : m}$	<i>Lex</i>
		$\frac{}{N \backslash (S / N) : \lambda x_1. \lambda x_2. ((A x_2) x_1)}$	<i>R3</i>	$\frac{}{(S \backslash N) / S : \lambda x_0. (x_0 m)}$	<i>R4</i>
		$N \backslash S : \lambda x_3. (\lambda x_0. (x_0 m) (\lambda x_1. \lambda x_2. ((A x_2) x_1) x_3))$			<i>R2</i>
		$=_{\beta-red.} \lambda x_3. ((\lambda x_1. \lambda x_2. ((A x_2) x_1) x_3) m)$			
		$=_{\beta-red.} \lambda x_3. (\lambda x_2. ((A x_2) x_3) m)$			
		$=_{\beta-red.} \lambda x_3. ((A m) x_3)$			
		$\frac{}{S : (\lambda x_3. ((A m) x_3) p)}$			<i>R1</i>

```

1.
s : ((A m) p), por redução-beta
  s : (lambda-x_3.((A m) x_3) p), por r1
    n : p, por lx
      pedro
    n\s : lambda-x_3.((A m) x_3), por redução-beta
      n\s : lambda-x_3.(lambda-x_2.((A x_2) x_3) m), por redução-beta
        n\s : lambda-x_3.((lambda-x_1.lambda-x_2.((A x_2) x_1) x_3) m), por redução-beta
          n\s : lambda-x_3.(lambda-x_0.(x_0 m) (lambda-x_1.lambda-x_2.((A x_2) x_1) x_3)), por r2
            n\s/n : lambda-x_1.lambda-x_2.((A x_2) x_1), por r3
              (n\s)/n : A, por lx
                ama
              (s/n)\s : lambda-x_0.(x_0 m), por r4
                n : m, por lx
                  maria

```

Análise: ajuda.

Pedro		ama		Maria	
$\frac{}{N : p}$	<i>Lex</i>	$\frac{}{(N \backslash S) / N : A}$	<i>Lex</i>	$\frac{}{N : m}$	<i>Lex</i>
		$\frac{}{N \backslash (S / N) : \lambda x_1. \lambda x_2. ((A x_2) x_1)}$	<i>R3</i>	$\frac{}{(S \backslash N) / S : \lambda x_0. (x_0 m)}$	<i>R4</i>
		$\frac{}{N \backslash S : \lambda x_3. (\lambda x_0. (x_0 m) (\lambda x_1. \lambda x_2. ((A x_2) x_1) x_3))}$			<i>R2</i>
		$=_{\beta-red.} \lambda x_3. ((\lambda x_1. \lambda x_2. ((A x_2) x_1) x_3) m)$			
		$=_{\beta-red.} \lambda x_3. (\lambda x_2. ((A x_2) x_3) m)$			
		$=_{\beta-red.} \lambda x_3. ((A m) x_3)$			
		$\frac{}{S : (\lambda x_3. ((A m) x_3) p)}$			<i>R1</i>
		$=_{\beta-red.} ((A m) p)$			

r1(N) ou r1 N	: R1 a partir de N
r2(N) ou r2 N	: R2 a partir de N
r3(N) ou r3 N	: R3 em N
r4(N) ou r4 N	: R4 em N (o assistente vai solicitar o conectivo principal)
r4(N, C)	: R4 em N usando C (/ ou \\\) como conectivo principal
red(N) ou red N	: redução-beta da representação semântica em N
desfaz(N) ou desfaz N	: desfaz a última operação em N
análise	: mostra novamente a análise já feita
nova	: solicita uma nova expressão para ser analisada
encerra	: encerra o assistente de análise gramatical

Análise: encerra.

Análise encerrada.

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5 Conclusions

- Mainly a didactic tool
- But a step by step processing model – psycholinguistics interest
- Easily adaptable to Classical Categorical Grammar and Lambek calculus
- To do list
 - complete lambda calculus?
 - linguistic data
 - syntactic variable presentation
 - graphical interface

References

- [1] Joel M. Cohen. The equivalence of two concepts of categorial grammar. *Information and Control*, 10:475–484, 1967.
- [2] Mary McGee Wood. *Categorial Grammars*. Routledge, London, 1993.

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