

Gramática Categorical: Introdução de $\backslash$ e de $/$

Luiz Arthur Pagani (UFPR)

1 Introdução de $\backslash$

- regra:

$$\boxed{
 \begin{array}{c}
 [Y : v]^n \qquad \vdots \\
 \vdots \qquad \vdots \\
 \hline
 X : \phi \\
 \hline
 Y \backslash X : \lambda v. \phi \quad I \backslash^n
 \end{array}
 }$$

- condições:
 - $v \in Var_{T_Y}$
 - v nova na derivação

[Baseado em [1, p. 156]]

2 Introdução de $/$

- regra:

$$\boxed{
 \begin{array}{c}
 \vdots \qquad [Y : v]^n \\
 \vdots \qquad \vdots \\
 \hline
 X : \phi \\
 \hline
 X/Y : \lambda v. \phi \quad I/n
 \end{array}
 }$$

- condições:
 - $v \in Var_{T_Y}$
 - v nova na derivação

[Baseado em [1, p. 156]]

3 Exemplos

3.1 “Pedro ama e Maria odeia Grosbilda”

3.1.1 “Pedro ama”

- inserção lexical:

$$\boxed{\frac{\text{Pedro}}{N : p}^L \quad \frac{\text{ama}}{(N \backslash S) / N : A}^L}$$

- suposição inicial:

$$\boxed{\frac{\text{Pedro}}{N : p}^L \quad \frac{\text{ama}}{(N \backslash S) / N : A}^L \quad [N : x]}$$

- eliminação de $/$:

$$\begin{array}{ccccccc}
 \underbrace{(N \backslash S)} & / & N & : & A & & N & : & x \\
 \downarrow & | & \downarrow & | & \downarrow & & \downarrow & | & \downarrow \\
 \boxed{X} & / & Y & : & \phi & + & Y & : & \alpha \Rightarrow X & : & (\phi \alpha) \\
 & & & & & & & & \downarrow & | & | & \downarrow & \downarrow & | \\
 & & & & & & & & \underbrace{N \backslash S} & : & (A \ x)
 \end{array}$$

- permite desenhar:

$$\boxed{
 \begin{array}{c}
 \text{Pedro} \\
 \hline
 N : p \quad L
 \end{array}
 \quad
 \begin{array}{c}
 \text{ama} \\
 \hline
 (N \backslash S) / N : A \quad L
 \end{array}
 \quad
 [N : x] \\
 \hline
 N \backslash S : (A \ x) \quad E/$$

- eliminação de \:

$$\begin{array}{ccccccc}
 N & : & p & & N & \backslash & S : \underbrace{(A x)} \\
 \downarrow & | & \downarrow & & \downarrow & | & \downarrow & | & \downarrow \\
 \boxed{Y : \alpha + Y \backslash X : \phi} & \Rightarrow & X : (& \phi & \alpha &) & \\
 & & \downarrow & | & | & \downarrow & \downarrow & | \\
 & & S : (& \underbrace{(A x)} & p &) &
 \end{array}$$

- permite agora desenhar:

$$\boxed{
 \begin{array}{c}
 \text{Pedro} \\
 \hline N : p \quad L \\
 \text{ama} \\
 \hline (N \backslash S) / N : A \quad L \quad [N : x]^1 \\
 \hline N \backslash S : (A x) \quad E/ \\
 \hline S : ((A x) p) \quad E\backslash
 \end{array}
 }$$

- introdução de $/$:

$$\begin{array}{c}
 \text{Pedro} \\
 \hline
 N : p \quad L
 \end{array}
 \quad
 \begin{array}{c}
 \text{ama} \\
 \hline
 (N \backslash S) / N : A \quad L
 \end{array}
 \quad
 [N : x]^1$$

$$\frac{\quad}{N \backslash S : (A \ x)} E/$$

$$\frac{\quad}{S : ((A \ x) \ p)} E\backslash$$

$$\frac{\quad}{S/N : \lambda x.((A \ x) \ p)} I/^1$$

3.1.2 “Maria odeia”

- inserção lexical:

$$\boxed{\frac{\text{Maria}}{N : m}^L \quad \frac{\text{odeia}}{(N \backslash S) / N : O}^L}$$

- suposição inicial:

$$\boxed{\frac{\text{Maria}}{N : m}^L \quad \frac{\text{odeia}}{(N \backslash S) / N : O}^L \quad [N : y]}$$

- eliminação de $/$:

$$\begin{array}{ccccccc}
 \underbrace{(N \backslash S)} & / & N & : & O & & N & : & y \\
 \downarrow & | & \downarrow & | & \downarrow & & \downarrow & | & \downarrow \\
 \boxed{X & / & Y & : & \phi & + & Y & : & \alpha} \Rightarrow X & : & (& \phi & \alpha &) \\
 & & & & & & \downarrow & | & | & \downarrow & \downarrow & | \\
 & & & & & & \underbrace{N \backslash S} & : & (& O & y &)
 \end{array}$$

- permite desenhar:

$$\boxed{
 \begin{array}{c}
 \text{Maria} \\
 \hline
 N : m \quad L
 \end{array}
 \quad
 \begin{array}{c}
 \text{odeia} \\
 \hline
 (N \backslash S) / N : O \quad L
 \end{array}
 \quad
 [N : y] \\
 \hline
 N \backslash S : (O \ y) \quad E/$$

- eliminação de \:

$$\begin{array}{ccccccc}
 N & : & m & & N & \backslash & S & : & \underbrace{(O \ y)} \\
 \downarrow & | & \downarrow & & \downarrow & | & \downarrow & | & \downarrow \\
 \boxed{Y & : & \alpha & + & Y & \backslash & X & : & \phi} & \Rightarrow & X & : & (& \phi & \alpha &) \\
 & & & & & & & & & & \downarrow & | & | & \downarrow & \downarrow & | \\
 & & & & & & & & & & S & : & (& \underbrace{(O \ y)} & m &)
 \end{array}$$

- permite agora desenhar:

$$\boxed{
 \begin{array}{c}
 \text{Maria} \\
 \hline
 N : m \quad L \\
 \hline
 \text{odeia} \\
 \hline
 (N \backslash S) / N : O \quad L \quad [N : y] \\
 \hline
 N \backslash S : (O \ y) \quad E/ \\
 \hline
 S : ((O \ y) \ m) \quad E\backslash
 \end{array}
 }$$

- introdução de /:

$$\begin{array}{c}
 \text{Maria} \\
 \hline N : m \quad L
 \end{array}
 \quad
 \begin{array}{c}
 \text{odeia} \\
 \hline (N \backslash S) / N : O \quad L
 \end{array}
 \quad
 [N : y]^2$$

$$\begin{array}{c}
 \hline N \backslash S : (O \ y) \quad E/
 \end{array}$$

$$\begin{array}{c}
 \hline S : ((O \ y) \ m) \quad E\backslash
 \end{array}$$

$$\begin{array}{c}
 \hline S / N : \lambda y.((O \ y) \ m) \quad I/{}^2
 \end{array}$$

3.1.3 “Pedro ama e Maria odeia Grosbilda”

- $\langle e, ((S/N) \backslash (S/N)) / (S/N), \lambda P. \lambda Q. \lambda z. ((Q z) \wedge (P z)) \rangle \in Lex$
- início da análise:

$\frac{\text{Pedro ama}}{S/N : \lambda x. ((A x) p)} D$	$\frac{e}{((S/N) \backslash (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z. ((Q z) \wedge (P z))} L$	$\frac{\text{Maria odeia}}{S/N : \lambda y. ((O y) m)} D$	$\frac{\text{Grosbilda}}{N : g} L$
---	--	---	------------------------------------

- eliminação de /:

$$\begin{array}{ccccccc}
 \underbrace{((S/N) \backslash (S/N))} & / & \underbrace{(S/N)} & : & \underbrace{\lambda P. \lambda Q. \lambda z. ((Q \ z) \wedge (P \ z))} & & \underbrace{S/N} : \underbrace{\lambda y. ((O \ y) \ m)} \\
 \downarrow & | & \downarrow & | & \downarrow & & \downarrow \quad | \quad \downarrow \\
 \boxed{X \quad / \quad Y \quad : \quad \phi \quad + \quad Y \quad : \quad \alpha}
 \end{array}$$

$$\Downarrow$$

$$\begin{array}{ccccccc}
 \boxed{X \quad : \quad (\quad \phi \quad \alpha \quad)} \\
 \downarrow & | & | & \downarrow & \downarrow & & | \\
 \underbrace{(S/N) \backslash (S/N)} & : & (& \underbrace{\lambda P. \lambda Q. \lambda z. ((Q \ z) \wedge (P \ z))} & \underbrace{\lambda y. ((O \ y) \ m)} &)
 \end{array}$$

- inclusão na análise:

$$\boxed{
 \begin{array}{c}
 \frac{\text{Pedro ama}}{S/N : \lambda x.((A\ x)\ p)} D \quad \frac{\frac{e}{((S/N) \backslash (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))} L \quad \frac{\text{Maria odeia}}{S/N : \lambda y.((O\ y)\ m)} D \quad \frac{\text{Grosbilda}}{N : g} L \\
 \hline
 (S/N) \backslash (S/N) : (\lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))\ \lambda y.((O\ y)\ m)) E/
 \end{array}
 }$$

- redução- β :

$$\begin{array}{c}
 (\quad \lambda \quad P \quad . \quad \underbrace{\lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))}_{\varphi} \quad \lambda y. ((O\ y)\ m) \quad) \\
 \begin{array}{ccccccc}
 | & | & \downarrow & | & & \downarrow & & \downarrow & | \\
 \hline
 (& \lambda & v & . & & \varphi & & \alpha &)
 \end{array} \\
 \Downarrow \\
 \boxed{\varphi[v \mapsto \alpha]} \\
 \downarrow \\
 \underbrace{\lambda Q. \lambda z. ((Q\ z) \wedge (\lambda y. ((O\ y)\ m)\ z))}
 \end{array}$$

- inclusão na análise:

$\frac{\text{Pedro ama}}{S/N : \lambda x.((A\ x)\ p)} D$	$\frac{e}{((S/N) \setminus (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))} L$	$\frac{\text{Maria odeia}}{S/N : \lambda y.((O\ y)\ m)} D$	$\frac{\text{Grosbilda}}{N : g} L$
$\frac{\lambda y.((O\ y)\ m)}{E/}$			
$\frac{(S/N) \setminus (S/N) : (\lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))\ \lambda y.((O\ y)\ m))}{=_{red.\beta} \lambda Q. \lambda z. ((Q\ z) \wedge (\lambda y.((O\ y)\ m)\ z))}$			

- redução- β em subfórmula:

$$\lambda Q. \lambda z. ((Q\ z) \wedge \underbrace{(\lambda y. ((O\ y)\ m)\ z)}_{\text{aqui}})$$

- 2a. redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & y & . & \underbrace{((O\ y)\ m)} & z &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{ & \lambda & v & . & \varphi & \alpha & }
 \end{array}$$

$$\Downarrow$$

$$\boxed{\varphi[v \mapsto \alpha]}$$

$$\downarrow$$

$$\underbrace{((O\ z)\ m)}$$

- resultado:

$$\underbrace{\lambda Q. \lambda z. ((Q\ z) \wedge \underbrace{(\lambda y. ((O\ y)\ m)\ z})}_{((O\ z)\ m)}$$

$$\lambda Q. \lambda z. ((Q\ z) \wedge ((O\ z)\ m))$$

- inclusão na análise:

$$\begin{array}{c}
 \frac{\text{Pedro ama}}{S/N : \lambda x.((A x) p)} D \quad \frac{\frac{e}{((S/N) \backslash (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z.((Q z) \wedge (P z))} L \quad \frac{\text{Maria odeia}}{S/N : \lambda y.((O y) m)} D \quad \frac{\text{Grosbilda}}{N : g} L}{\frac{(S/N) \backslash (S/N) : (\lambda P. \lambda Q. \lambda z.((Q z) \wedge (P z)) \lambda y.((O y) m))}{=_{red.\beta} \lambda Q. \lambda z.((Q z) \wedge (\lambda y.((O y) m) z))}{=_{red.\beta} \lambda Q. \lambda z.((Q z) \wedge ((O z) m))} E/
 \end{array}$$

- eliminação de \:

$$\begin{array}{c}
 \begin{array}{ccccccc}
 S/N & : & \underbrace{\lambda x.((A x) p)} & & (S/N) & \backslash & (S/N) : \underbrace{\lambda Q. \lambda z.((Q z) \wedge ((O z) m))} \\
 \downarrow & | & \downarrow & & \downarrow & | & \downarrow \\
 Y & : & \alpha & + & Y & \backslash & X : \phi
 \end{array} \\
 \Downarrow \\
 \begin{array}{c}
 \boxed{X : (\quad \phi \quad \alpha \quad)} \\
 \downarrow \quad | \quad | \quad \downarrow \quad \downarrow \quad | \\
 S/N : (\quad \underbrace{\lambda Q. \lambda z.((Q z) \wedge ((O z) m))} \quad \underbrace{\lambda x.((A x) p)} \quad)
 \end{array}
 \end{array}$$

- inclusão na análise:

$$\begin{array}{c}
 \frac{\text{Pedro ama}}{S/N : \lambda x.((A\ x)\ p)} D \quad \frac{\frac{e}{((S/N) \setminus (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z.((Q\ z) \wedge (P\ z))} L \quad \frac{\text{Maria odeia}}{S/N : \lambda y.((O\ y)\ m)} D \quad \frac{\text{Grosbilda}}{N : g} L}{\frac{(S/N) \setminus (S/N) : (\lambda P. \lambda Q. \lambda z.((Q\ z) \wedge (P\ z))\ \lambda y.((O\ y)\ m))}{=_{red.\beta} \lambda Q. \lambda z.((Q\ z) \wedge (\lambda y.((O\ y)\ m)\ z))}{=_{red.\beta} \lambda Q. \lambda z.((Q\ z) \wedge ((O\ z)\ m))} E/} \\
 \hline
 S/N : (\lambda Q. \lambda z.((Q\ z) \wedge ((O\ z)\ m))\ \lambda x.((A\ x)\ p)) E\backslash
 \end{array}$$

- redução- β :

$$\begin{array}{c}
 (\quad \lambda \quad Q \quad . \quad \underbrace{\lambda z.((Q\ z) \wedge ((O\ z)\ m))}_{\varphi} \quad \underbrace{\lambda x.((A\ x)\ p)}_{\alpha} \quad) \\
 \begin{array}{ccccccc}
 | & | & \downarrow & | & & \downarrow & | \\
 \hline
 (& \lambda & v & . & \varphi & & \alpha &)
 \end{array} \\
 \Downarrow \\
 \boxed{\varphi[v \mapsto \alpha]} \\
 \downarrow \\
 \overbrace{\lambda z.((\lambda x.((A\ x)\ p)\ z) \wedge ((O\ z)\ m))}
 \end{array}$$

- inclusão na análise:

$$\begin{array}{c}
 \frac{\text{Pedro ama}}{S/N : \lambda x.((A\ x)\ p)} D \quad \frac{\frac{e}{((S/N) \setminus (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))} L \quad \frac{\text{Maria odeia}}{S/N : \lambda y.((O\ y)\ m)} D \quad \frac{\text{Grosbilda}}{N : g} L}{\frac{(\lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))\ \lambda y.((O\ y)\ m))}{=_{red.\beta} \lambda Q. \lambda z. ((Q\ z) \wedge (\lambda y.((O\ y)\ m)\ z))} E/ \quad \frac{=_{red.\beta} \lambda Q. \lambda z. ((Q\ z) \wedge ((O\ z)\ m))}{S/N : (\lambda Q. \lambda z. ((Q\ z) \wedge ((O\ z)\ m))\ \lambda x.((A\ x)\ p))} E\backslash \\
 \frac{=_{red.\beta} \lambda z. ((\lambda x.((A\ x)\ p)\ z) \wedge ((O\ z)\ m))}{S/N : (\lambda z. ((\lambda x.((A\ x)\ p)\ z) \wedge ((O\ z)\ m))} E\backslash
 \end{array}$$

- outra redução- β em subfórmula:

$$\lambda z. (\underbrace{(\lambda x.((A\ x)\ p)\ z)}_{\text{aqui}} \wedge ((O\ z)\ m))$$

- nova redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & x & . & \underbrace{((A\ x)\ p)} & z &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{(& \lambda & v & . & \varphi & \alpha &)} \\
 & & & & \Downarrow & & \\
 & & & & \boxed{\varphi[v \mapsto \alpha]} & & \\
 & & & & \downarrow & & \\
 & & & & \underbrace{((A\ z)\ p)} & &
 \end{array}$$

- resultado:

$$\begin{array}{c}
 \lambda z. (\underbrace{(\lambda x. ((A\ x)\ p)\ z)}_{((A\ z)\ p)} \wedge ((O\ z)\ m)) \\
 \underbrace{\hspace{10em}} \\
 \lambda z. ((A\ z)\ p) \wedge ((o\ z)\ m))
 \end{array}$$

- inclusão na análise:

$$\begin{array}{c}
 \frac{\text{Pedro ama}}{S/N : \lambda x.((A x) p)} D \quad \frac{\frac{e}{((S/N) \setminus (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z.((Q z) \wedge (P z))} L \quad \frac{\text{Maria odeia}}{S/N : \lambda y.((O y) m)} D \quad \frac{\text{Grosbilda}}{N : g} L}{\frac{\frac{\frac{(S/N) \setminus (S/N) : (\lambda P. \lambda Q. \lambda z.((Q z) \wedge (P z)) \lambda y.((O y) m))}{=_{red.\beta} \lambda Q. \lambda z.((Q z) \wedge (\lambda y.((O y) m) z))}{=_{red.\beta} \lambda Q. \lambda z.((Q z) \wedge ((O z) m))}}{S/N : (\lambda Q. \lambda z.((Q z) \wedge ((O z) m)) \lambda x.((A x) p))}{E/} \\
 \frac{=_{red.\beta} \lambda z.((\lambda x.((A x) p) z) \wedge ((O z) m))}{=_{red.\beta} \lambda z.(((A z) p) \wedge ((O z) m))} E\backslash
 \end{array}$$

- última eliminação de /:

$$\begin{array}{c}
 \begin{array}{ccccc}
 S & / & N & : & \underbrace{\lambda z.(((A z) p) \wedge ((O z) m))} \\
 \downarrow & | & \downarrow & | & \downarrow \\
 X & / & Y & : & \phi
 \end{array}
 \quad
 \begin{array}{ccc}
 N & : & g \\
 \downarrow & | & \downarrow \\
 + & Y & : \alpha
 \end{array} \\
 \Downarrow \\
 \begin{array}{ccccc}
 X & : & (& \phi & \alpha &) \\
 \downarrow & | & | & \downarrow & \downarrow & | \\
 S & : & (& \underbrace{\lambda z.(((A z) p) \wedge ((O z) m))} & g &)
 \end{array}
 \end{array}$$

- inclusão na análise:

$$\begin{array}{c}
 \frac{\text{Pedro ama}}{S/N : \lambda x.((A\ x)\ p)} D \quad \frac{e}{((S/N) \setminus (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z))} L \quad \frac{\text{Maria odeia}}{S/N : \lambda y.((O\ y)\ m)} D \quad \frac{\text{Grosbilda}}{N : g} L \\
 \hline
 \frac{\lambda P. \lambda Q. \lambda z. ((Q\ z) \wedge (P\ z)) \quad \lambda y.((O\ y)\ m)}{=_{red.\beta} \lambda Q. \lambda z. ((Q\ z) \wedge (\lambda y.((O\ y)\ m)\ z))} E/ \\
 \hline
 \frac{=_{red.\beta} \lambda Q. \lambda z. ((Q\ z) \wedge ((O\ z)\ m))}{S/N : (\lambda Q. \lambda z. ((Q\ z) \wedge ((O\ z)\ m)) \quad \lambda x.((A\ x)\ p))} E\backslash \\
 \hline
 \frac{=_{red.\beta} \lambda z. ((\lambda x.((A\ x)\ p)\ z) \wedge ((O\ z)\ m))}{=_{red.\beta} \lambda z. (((A\ z)\ p) \wedge ((O\ z)\ m))} \\
 \hline
 S : (\lambda z. (((A\ z)\ p) \wedge ((O\ z)\ m))\ g) E/
 \end{array}$$

- última redução- β :

$$\begin{array}{c}
 (\quad \lambda \quad z \quad . \quad \underbrace{(((A\ z)\ p) \wedge ((O\ z)\ m))} \quad g \quad) \\
 \begin{array}{ccccccc}
 | & | & \downarrow & | & & \downarrow & \downarrow & | \\
 (& \lambda & v & . & & \varphi & & \alpha &)
 \end{array} \\
 \Downarrow \\
 \boxed{\varphi[v \mapsto \alpha]} \\
 \downarrow \\
 \overbrace{(((A\ g)\ p) \wedge ((O\ g)\ m))}
 \end{array}$$

- | | | | |
|--|--|--|--|
| $\frac{\text{Pedro ama}}{S/N : \lambda x.((A \ x) \ p)} \quad D$ | $\frac{e}{((S/N) \setminus (S/N)) / (S/N) : \lambda P. \lambda Q. \lambda z. ((Q \ z) \wedge (P \ z))} \quad L$ | $\frac{\text{Maria odeia}}{S/N : \lambda y.((O \ y) \ m)} \quad D$ | $\frac{\text{Grosbilda}}{N : g} \quad L$ |
| | $\frac{\lambda P. \lambda Q. \lambda z. ((Q \ z) \wedge (P \ z)) \quad \lambda y.((O \ y) \ m)}{E/}$ | | |
| | $\frac{\begin{aligned} &(S/N) \setminus (S/N) : \\ &(\lambda P. \lambda Q. \lambda z. ((Q \ z) \wedge (P \ z)) \quad \lambda y.((O \ y) \ m)) \\ &=_{red.\beta} \lambda Q. \lambda z. ((Q \ z) \wedge (\lambda y.((O \ y) \ m) \ z)) \\ &=_{red.\beta} \lambda Q. \lambda z. ((Q \ z) \wedge ((O \ z) \ m)) \end{aligned}}{E \setminus}$ | | |
| | $\frac{\begin{aligned} &S/N : (\lambda Q. \lambda z. ((Q \ z) \wedge ((O \ z) \ m)) \quad \lambda x.((A \ x) \ p)) \\ &=_{red.\beta} \lambda z. ((\lambda x.((A \ x) \ p) \ z) \wedge ((O \ z) \ m)) \\ &=_{red.\beta} \lambda z. (((A \ z) \ p) \wedge ((O \ z) \ m)) \end{aligned}}{E/}$ | | |
| | $\frac{\begin{aligned} &S : (\lambda z. (((A \ z) \ p) \wedge ((O \ z) \ m)) \ g) \\ &=_{red.\beta} (((A \ g) \ p) \wedge ((O \ g) \ m)) \end{aligned}}{E/}$ | | |

3.2 “todo homem ama uma mulher”

- sentença ambígua:
 - para todo homem há uma mulher tal que ele a ama
(uma mulher para cada homem)
 - há uma mulher tal que todo homem a ama
(uma única mulher para todos os homens)
- itens lexicais que faltam para a análise:
 - $\langle \text{homem}, N_c, H \rangle \in Lex$
 - $\langle \text{mulher}, N_c, M \rangle \in Lex$

3.2.1 “todo homem”

- inserção lexical:

$$\boxed{\frac{\text{todo}}{(S/(N \setminus S))/N_c : \lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x))}^L \quad \frac{\text{homem}}{N_c : H}^L}$$

- eliminação de /:

$$\boxed{\begin{array}{ccccccc} \underbrace{(S/(N \setminus S))} & / & N_c & : & \underbrace{\lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x))} & & N_c & : & H \\ \downarrow & | & \downarrow & | & \downarrow & & \downarrow & | & \downarrow \\ X & / & Y & : & \phi & + & Y & : & \alpha \end{array}}$$

$\Downarrow$

$$\boxed{\begin{array}{ccccccc} X & : & (& \phi & \alpha &) \\ \downarrow & | & | & \downarrow & \downarrow & | \\ \underbrace{S/(N \setminus S)} & : & (& \underbrace{\lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x))} & H &) \end{array}}$$

- incluindo a eliminação no diagrama:

$$\boxed{\frac{\frac{\text{todo}}{(S/(N \setminus S))/N_c : \lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x)))^L \quad \frac{\text{homem}}{N_c : H}^L}{S/(N \setminus S) : (\lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x)) \ H)}^{E/}}$$

- redução- β :

$$\begin{array}{c} \left(\begin{array}{cccc} \lambda & P & . & \underbrace{\lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x))} \\ | & | & \downarrow & | \\ \lambda & v & . & \varphi \end{array} \quad H \right) \\ \downarrow \\ \boxed{\varphi[v \mapsto \alpha]} \\ \downarrow \\ \underbrace{\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \end{array}$$

- incluindo a redução- β no diagrama:

$$\boxed{
 \begin{array}{c}
 \text{todo} \qquad \qquad \text{homem} \\
 \hline
 (S/(N \backslash S))/N_c : \lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x)) \quad \quad \quad N_c : H \\
 \hline
 S/(N \backslash S) : (\lambda P. \lambda Q. \forall x. ((P \ x) \rightarrow (Q \ x)) \ H) \\
 \qquad \qquad \qquad =_{red.\beta} \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))
 \end{array}
 }$$

- finalizado

3.2.2 “uma mulher”

- inserção lexical:

$$\boxed{\frac{\text{uma}}{((S/N)\backslash S)/N_c : \lambda P_1.\lambda Q_1.\exists y.((P_1 y) \wedge (Q_1 y))}^L \quad \frac{\text{mulher}}{N_c : M}^L}$$

- eliminação de /:

$$\boxed{\begin{array}{ccccccc} \underbrace{((S/N)\backslash S)} & / & N_c & : & \underbrace{\lambda P_1.\lambda Q_1.\exists y.((P_1 y) \wedge (Q_1 y))} & & N_c : M \\ \downarrow & | & \downarrow & | & \downarrow & & \downarrow | \downarrow \\ X & / & Y & : & \phi & + & Y : \alpha \end{array}}$$

$\Downarrow$

$$\boxed{\begin{array}{ccccccc} X & : & (& \phi & \alpha &) \\ \downarrow & | & | & \downarrow & \downarrow & | \\ \underbrace{(S/N)\backslash S} & : & (& \underbrace{\lambda P_1.\lambda Q_1.\exists y.((P_1 y) \wedge (Q_1 y))} & M &) \end{array}}$$

- incluindo a eliminação no diagrama:

$$\boxed{\frac{\frac{\text{uma}}{((S/N)\backslash S)/N_c : \lambda P_1.\lambda Q_1.\exists y.((P_1 y) \wedge (Q_1 y))}^L \quad \frac{\text{mulher}}{N_c : M}^L}{(S/N)\backslash S : (\lambda P_1.\lambda Q_1.\exists y.((P_1 y) \wedge (Q_1 y)) M)}^{E/}}$$

- redução- β :

$$\begin{array}{c} \left(\begin{array}{c} \lambda \quad P_1 \quad . \quad \underbrace{\lambda Q_1.\exists y.((P_1 y) \wedge (Q_1 y))}_{\text{}} \quad M \end{array} \right) \\ \begin{array}{ccccccc} | & | & \downarrow & | & & \downarrow & & \downarrow & | \\ \boxed{\left(\begin{array}{c} \lambda \quad v \quad . \quad \varphi \quad \alpha \end{array} \right)} & & & & & & & & \end{array} \\ \Downarrow \\ \boxed{\varphi[v \mapsto \alpha]} \\ \downarrow \\ \underbrace{\lambda Q_1.\exists y.((M y) \wedge (Q_1 y))}_{\text{}} \end{array}$$

- incluindo a redução- β no diagrama:

$$\boxed{
 \begin{array}{c}
 \text{uma} \qquad \qquad \qquad \text{mulher} \\
 \hline
 ((S/N) \backslash S) / N_c : \lambda P_1. \lambda Q_1. \exists y. ((P_1 \ y) \wedge (Q_1 \ y)) \quad L \quad N_c : M \quad L \\
 \hline
 (S/N) \backslash S : (\lambda P_1. \lambda Q_1. \exists y. ((P_1 \ y) \wedge (Q_1 \ y)) \ M) \quad E/ \\
 \qquad \qquad \qquad =_{red.\beta} \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))
 \end{array}
 }$$

- finalizado

3.2.3 Preparação de “ama”

- postulação de duas suposições:

$$\boxed{\begin{array}{c} \text{ama} \\ \hline [N : z]^1 \quad (N \backslash S) / N : A \quad L \quad [N : w]^2 \end{array}}$$

- eliminação de /:

$$\begin{array}{ccccccc} \underbrace{(N \backslash S)} & / & N & : & A & & N & : & w \\ \downarrow & | & \downarrow & | & \downarrow & & \downarrow & | & \downarrow \\ \boxed{X} & / & Y & : & \phi & + & Y & : & \alpha \Rightarrow X & : & (& \phi & \alpha &) \\ & & & & & & & & \downarrow & | & | & \downarrow & \downarrow & | \\ & & & & & & & & \underbrace{N \backslash S} & : & (& A & w &) \end{array}$$

- incluindo a eliminação de $/$ no diagrama:

$$\boxed{\begin{array}{c} \text{ama} \\ \frac{[N : z]^1 \quad \frac{\overline{(N \backslash S) / N : A}^L \quad [N : w]^2}{N \backslash S : (A w)} E/}{N \backslash S : (A w)} \end{array}}$$

- eliminação de $\backslash$:

$$\begin{array}{ccccccc} N & : & z & & N & \backslash & S & : & \underbrace{(A w)} \\ \downarrow & | & \downarrow & & \downarrow & | & \downarrow & | & \downarrow \\ \boxed{Y & : & \alpha & + & Y & \backslash & X & : & \phi} & \Rightarrow & X & : & (& \phi & \alpha &) \\ & & & & & & & & & & \downarrow & | & | & \downarrow & \downarrow & | \\ & & & & & & & & & & S & : & (& \underbrace{(A w)} & z &) \end{array}$$

- incluindo a eliminação de $\backslash$ no diagrama:

$$\boxed{
 \begin{array}{c}
 \text{ama} \\
 \hline
 [N : z]^1 \quad \frac{(N \backslash S) / N : A}{[N : w]^2} \quad L \\
 \hline
 N \backslash S : (A \ w) \\
 \hline
 S : ((A \ w) \ z) \quad E \backslash
 \end{array}
 }$$

- finalizado

3.2.4 Primeira interpretação

- descarregando a segunda hipótese — introdução de /:

$\frac{\text{todo homem}}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \quad D$	$\frac{\text{ama}}{S : ((A \ w) \ z)} \quad D$	$\frac{\text{uma mulher}}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} \quad D$
	$\frac{((A \ w) \ z)}{S/N : \lambda w. ((A \ w) \ z)} \quad I/2$	

- eliminação de \:

$$\begin{array}{c}
 \underbrace{S/N} : \underbrace{\lambda w.((A w) z)} \qquad \underbrace{(S/N)} \setminus S : \underbrace{\lambda Q_1.\exists y.((M y) \wedge (Q_1 y))} \\
 \downarrow \quad | \qquad \downarrow \qquad \downarrow \quad | \quad \downarrow \quad | \qquad \downarrow \\
 \boxed{Y : \alpha \quad + \quad Y \setminus X : \phi}
 \end{array}$$

$\Downarrow$

$$\begin{array}{c}
 \boxed{X : (\qquad \phi \qquad \qquad \alpha \qquad)} \\
 \downarrow \quad | \quad | \qquad \downarrow \qquad \downarrow \qquad | \\
 S : (\quad \underbrace{\lambda Q_1.\exists y.((M y) \wedge (Q_1 y))} \quad \underbrace{\lambda w.((A w) z)} \quad)
 \end{array}$$

- incluindo no diagrama a eliminação de \:

$$\begin{array}{c}
 \frac{\text{todo homem}}{S/(N \backslash S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} D \quad \frac{\text{ama}}{S : ((A \ w) \ z)} D \quad \frac{\text{uma mulher}}{(S/N) \backslash S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} D \\
 \frac{((A \ w) \ z)}{S/N : \lambda w. ((A \ w) \ z)} I/2 \\
 \frac{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. ((A \ w) \ z))}{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. ((A \ w) \ z))} E\backslash
 \end{array}$$

- redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & Q_1 & . & \underbrace{\exists y.((M \ y) \wedge (Q_1 \ y))} & \underbrace{\lambda w.((A \ w) \ z)} &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{ & \lambda & v & . & \varphi & \alpha & } \\
 & & & & \Downarrow & & \\
 & & & & \boxed{\varphi[v \mapsto \alpha]} & & \\
 & & & & \downarrow & & \\
 & & & & \underbrace{\exists y.((M \ y) \wedge (\lambda w.((A \ w) \ z) \ y))} & &
 \end{array}$$

- incluindo a redução- β no diagrama:

$$\begin{array}{c}
 \frac{\text{todo homem}}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))} \quad D \quad \frac{\text{ama}}{S : ((A\ w)\ z)} \quad D \quad \frac{\text{uma mulher}}{(S/N)\backslash S : \lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))} \quad D \\
 \frac{((A\ w)\ z)}{S/N : \lambda w.((A\ w)\ z)} \quad I/2 \\
 \frac{S : (\lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))\ \lambda w.((A\ w)\ z))}{\begin{array}{l} S : (\lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))\ \lambda w.((A\ w)\ z)) \\ =_{red.\beta} \exists y.((M\ y) \wedge (\lambda w.((A\ w)\ z)\ y)) \end{array}} \quad E\backslash
 \end{array}$$

- subfórmula a ser reduzida:

$$\exists y.((M\ y) \wedge \underbrace{(\lambda w.((A\ w)\ z)\ y)}_{\text{aqui}})$$

- redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & w & . & \underbrace{((A\ w)\ z)} & y &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{(& \lambda & v & . & \varphi & \alpha &) \Rightarrow \varphi[v \mapsto \alpha]} \\
 & & & & & \downarrow & \\
 & & & & & \underbrace{((A\ y)\ z)} &
 \end{array}$$

- resultado:

$$\begin{array}{c}
 \exists y.((M\ y) \wedge \underbrace{(\lambda w.((A\ w)\ z)\ y)}_{((A\ y)\ z)}) \\
 \underbrace{\hspace{10em}} \\
 \exists y.((M\ y) \wedge ((A\ y)\ z))
 \end{array}$$

- incluindo a segunda redução- β no diagrama:

$$\begin{array}{c}
 \frac{\text{todo homem}}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \quad D \quad \frac{\text{ama}}{S : ((A \ w) \ z)} \quad D \quad \frac{\text{uma mulher}}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} \quad D \\
 \frac{((A \ w) \ z)}{S/N : \lambda w. ((A \ w) \ z)} \quad I/2 \\
 \frac{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. ((A \ w) \ z))}{\begin{array}{l} =_{red.\beta} \exists y. ((M \ y) \wedge (\lambda w. ((A \ w) \ z) \ y)) \\ =_{red.\beta} \exists y. ((M \ y) \wedge ((A \ y) \ z)) \end{array}} \quad E \setminus
 \end{array}$$

- descarregando a primeira hipótese — introdução de \:

$$\begin{array}{c}
 \frac{\text{todo homem}}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \quad D \quad \frac{\text{ama}}{S : ((A \ w) \ z)} \quad D \quad \frac{\text{uma mulher}}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} \quad D \\
 \frac{((A \ w) \ z)}{S/N : \lambda w. ((A \ w) \ z)} \quad I/2 \\
 \frac{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. ((A \ w) \ z))}{=_{red.\beta} \exists y. ((M \ y) \wedge (\lambda w. ((A \ w) \ z) \ y))} \quad E \setminus \\
 \frac{=_{red.\beta} \exists y. ((M \ y) \wedge ((A \ y) \ z))}{N \setminus S : \lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))} \quad I \setminus 1
 \end{array}$$

- eliminação de /:

$$\begin{array}{c}
 \begin{array}{ccccccc}
 S & / & \underbrace{(N \backslash S)} & : & \underbrace{\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} & N \backslash S & : & \underbrace{\lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))} \\
 \downarrow & | & \downarrow & | & \downarrow & \downarrow & | & \downarrow \\
 \boxed{X \quad / \quad Y \quad : \quad \phi \quad + \quad Y \quad : \quad \alpha}
 \end{array} \\
 \Downarrow \\
 \begin{array}{ccccccc}
 \boxed{X \quad : \quad (\quad \phi \quad \alpha \quad)} \\
 \downarrow \quad | \quad | \quad \downarrow \quad \downarrow \quad \downarrow \quad | \\
 S \quad : \quad (\quad \underbrace{\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \quad \underbrace{\lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))} \quad)
 \end{array}
 \end{array}$$

- incluindo no diagrama a eliminação de /:

$\frac{\text{todo homem}}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))} \quad D$	$\frac{\text{ama}}{S : ((A\ w)\ z)} \quad D$	$\frac{\text{uma mulher}}{(S/N)\backslash S : \lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))} \quad D$
	$\frac{((A\ w)\ z)}{\lambda w.((A\ w)\ z)} \quad I/2$	
	$\frac{S/N : \lambda w.((A\ w)\ z)}{S : (\lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))\ \lambda w.((A\ w)\ z))} \quad E\backslash$	
	$\begin{aligned} &=_{red.\beta} \exists y.((M\ y) \wedge (\lambda w.((A\ w)\ z)\ y)) \\ &=_{red.\beta} \exists y.((M\ y) \wedge ((A\ y)\ z)) \end{aligned}$	
	$\frac{\lambda z.\exists y.((M\ y) \wedge ((A\ y)\ z))}{N\backslash S : \lambda z.\exists y.((M\ y) \wedge ((A\ y)\ z))} \quad I\backslash^1$	
	$\frac{S : (\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))\ \lambda z.\exists y.((M\ y) \wedge ((A\ y)\ z)))}{S : (\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))\ \lambda z.\exists y.((M\ y) \wedge ((A\ y)\ z)))} \quad E\backslash$	

- redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & Q & . & \underbrace{\forall x.((H\ x) \rightarrow (Q\ x))} & \underbrace{\lambda z.\exists y.((M\ y) \wedge ((A\ y)\ z))} &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{(& \lambda & v & . & \varphi & \alpha &)}
 \end{array}$$

$$\Downarrow$$

$$\boxed{\varphi[v \mapsto \alpha]}$$

$$\downarrow$$

$$\overbrace{\forall x.((H\ x) \rightarrow (\lambda z.\exists y.((M\ y) \wedge ((A\ y)\ z))\ x))}$$

- incluindo no diagrama a redução- β :

todo homem	ama	uma mulher
$\frac{}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} D$	$\frac{}{S : ((A \ w) \ z)} D$	$\frac{}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} D$
	$\frac{}{S/N : \lambda w. ((A \ w) \ z)} I/2$	
	$\frac{}{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. ((A \ w) \ z))} E \setminus$	
	$\frac{}{=_{red.\beta} \exists y. ((M \ y) \wedge (\lambda w. ((A \ w) \ z) \ y))}$	
	$\frac{}{=_{red.\beta} \exists y. ((M \ y) \wedge ((A \ y) \ z))} I \setminus^1$	
	$\frac{}{N \setminus S : \lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))} E \setminus$	
	$\frac{}{S : (\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) \ \lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z)))}$	
	$\frac{}{=_{red.\beta} \forall x. ((H \ x) \rightarrow (\lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z)) \ x))}$	

- subfórmula a ser reduzida:

$$\forall x.((H \ x) \rightarrow \underbrace{(\lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))) \ x}_{\text{aqui}})$$

- redução- β :

$$\begin{array}{ccccccc} (& \lambda & z & . & \underbrace{\exists y. ((M \ y) \wedge ((A \ y) \ z))} & x &) \\ | & | & \downarrow & | & \downarrow & \downarrow & | \\ \boxed{(& \lambda & v & . & \varphi & \alpha &) \Rightarrow \varphi[v \mapsto \alpha]} & & & & & \\ & & & & & & \downarrow \\ & & & & & & \underbrace{\exists y. ((M \ y) \wedge ((A \ y) \ x))} \end{array}$$

- resultado:

$$\forall x.((H \ x) \rightarrow \underbrace{(\lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))) \ x}_{\exists y. ((M \ y) \wedge ((A \ y) \ x))})$$

$$\underbrace{\forall x.((H \ x) \rightarrow \exists y. ((M \ y) \wedge ((A \ y) \ x)))}$$

- incluindo no diagrama a segunda redução- β :

todo homem	ama	uma mulher
$\frac{}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} D$	$\frac{}{S : ((A \ w) \ z)} D$	$\frac{}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} D$
	$\frac{}{S/N : \lambda w. ((A \ w) \ z)} I/2$	
	$\frac{}{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. ((A \ w) \ z))} E \setminus$	
	$\frac{}{=_{red.\beta} \exists y. ((M \ y) \wedge (\lambda w. ((A \ w) \ z) \ y))}$	
	$\frac{}{=_{red.\beta} \exists y. ((M \ y) \wedge ((A \ y) \ z))}$	
	$\frac{}{N \setminus S : \lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z))} I \setminus 1$	
	$\frac{}{S : (\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) \ \lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z)))} E \setminus$	
	$\frac{}{=_{red.\beta} \forall x. ((H \ x) \rightarrow (\lambda z. \exists y. ((M \ y) \wedge ((A \ y) \ z)) \ x))}$	
	$\frac{}{=_{red.\beta} \forall x. ((H \ x) \rightarrow \exists y. ((M \ y) \wedge ((A \ y) \ x)))}$	

3.2.5 Segunda interpretação

- descarregando a primeira hipótese — introdução de \:

$\frac{\text{todo homem}}{S/(N \backslash S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \quad D$	$\frac{\text{ama}}{S : ((A \ w) \ z)} \quad D$	$\frac{\text{uma mulher}}{(S/N) \backslash S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} \quad D$
	$\frac{((A \ w) \ z)}{N \backslash S : \lambda z. ((A \ w) \ z)} \quad I \backslash^1$	

- eliminação de /:

$$\begin{array}{ccccccc}
 S & / & \underbrace{(N \backslash S)} & : & \underbrace{\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} & & \underbrace{N \backslash S} : \underbrace{\lambda z. ((A \ w) \ z)} \\
 \downarrow & | & \downarrow & | & \downarrow & & \downarrow \quad | \quad \downarrow \\
 \boxed{X \quad / \quad Y \quad : \quad \phi \quad + \quad Y \quad : \quad \alpha} \\
 \\
 \Downarrow \\
 \boxed{X \quad : \quad (\quad \phi \quad \alpha \quad)} \\
 \downarrow \quad | \quad | \quad \downarrow \quad \downarrow \quad | \\
 S \quad : \quad (\quad \underbrace{\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} \quad \underbrace{\lambda z. ((A \ w) \ z)} \quad)
 \end{array}$$

- incluindo no diagrama a eliminação de $/$:

$$\begin{array}{c}
 \begin{array}{ccc}
 \text{todo homem} & \text{ama} & \text{uma mulher} \\
 \hline
 S/(N \backslash S) : & S : & (S/N) \backslash S : \\
 \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) & ((A \ w) \ z) & \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))
 \end{array} \\
 \begin{array}{ccc}
 & & \\
 & \hline
 & N \backslash S : & \\
 & \lambda z. ((A \ w) \ z) & \\
 & \hline
 & & E/
 \end{array} \\
 \hline
 S : (\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) \ \lambda z. ((A \ w) \ z))
 \end{array}$$

- redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & Q & . & \underbrace{\forall x.((H\ x) \rightarrow (Q\ x))} & \underbrace{\lambda z.((A\ w)\ z)} &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{ & \lambda & v & . & \varphi & \alpha & } \\
 & & & & \Downarrow & & \\
 & & & & \boxed{\varphi[v \mapsto \alpha]} & & \\
 & & & & \downarrow & & \\
 & & & & \underbrace{\forall x.((H\ x) \rightarrow (\lambda z.((A\ w)\ z)\ x))} & &
 \end{array}$$

- incluindo a redução- β no diagrama:

$$\begin{array}{c}
 \frac{\text{todo homem}}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))} \quad D \quad \frac{\text{ama}}{S : ((A\ w)\ z)} \quad D \quad \frac{\text{uma mulher}}{\lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))} \quad D \\
 \frac{\quad}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))} \quad S/(N \setminus S) : \\
 \frac{\quad}{\lambda z.((A\ w)\ z)} \quad N \setminus S : \\
 \frac{\quad}{S : (\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))\ \lambda z.((A\ w)\ z))} \quad I \setminus^1 \\
 \frac{\quad}{\quad} \quad E/ \\
 \frac{\quad}{\quad} \quad =_{red.\beta} \forall x.((H\ x) \rightarrow (\lambda z.((A\ w)\ z)\ x))
 \end{array}$$

- subfórmula a ser reduzida:

$$\forall x.((H\ x) \rightarrow \underbrace{(\lambda z.((A\ w)\ z)\ x)}_{\text{aqui}})$$

- redução- β :

$$\begin{array}{ccccccc}
 (& \lambda & z & . & \underbrace{((A\ w)\ z)} & x &) \\
 | & | & \downarrow & | & \downarrow & \downarrow & | \\
 \boxed{(& \lambda & v & . & \varphi & \alpha &) \Rightarrow \varphi[v \mapsto \alpha]} \\
 & & & & & \downarrow & \\
 & & & & & \underbrace{((A\ w)\ x))} &
 \end{array}$$

- resultado:

$$\begin{array}{c}
 \forall x.((H\ x) \rightarrow \underbrace{(\lambda z.((A\ w)\ z)\ x)}_{((A\ w)\ x)}) \\
 \underbrace{\hspace{10em}} \\
 \forall x.((H\ x) \rightarrow ((A\ w)\ x))
 \end{array}$$

- incluindo a segunda redução- β no diagrama:

$$\begin{array}{c}
 \frac{\text{todo homem}}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))} \quad D \quad \frac{\text{ama}}{S : ((A\ w)\ z)} \quad D \quad \frac{\text{uma mulher}}{(S/N)\backslash S : \lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))} \quad D \\
 \frac{\quad}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x)) \quad ((A\ w)\ z)} \quad I\backslash^1 \\
 \frac{\quad}{\lambda z.((A\ w)\ z)} \quad N\backslash S : \\
 \frac{\quad}{S : (\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x)) \lambda z.((A\ w)\ z))} \quad E/ \\
 \begin{array}{l}
 =_{red.\beta} \forall x.((H\ x) \rightarrow (\lambda z.((A\ w)\ z)\ x)) \\
 =_{red.\beta} \forall x.((H\ x) \rightarrow ((A\ w)\ x))
 \end{array}
 \end{array}$$

- descarregando a segunda hipótese — introdução de /:

$$\begin{array}{c}
 \frac{\text{todo homem}}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} D \quad \frac{\text{ama}}{S : ((A \ w) \ z)} D \quad \frac{\text{uma mulher}}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} D \\
 \frac{\quad}{\lambda z. ((A \ w) \ z)} I \setminus^1 \\
 \frac{\quad}{S : (\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) \ \lambda z. ((A \ w) \ z))} E/ \\
 \frac{\quad}{\begin{array}{l} =_{red.\beta} \forall x. ((H \ x) \rightarrow (\lambda z. ((A \ w) \ z) \ x)) \\ =_{red.\beta} \forall x. ((H \ x) \rightarrow ((A \ w) \ x)) \end{array}} I/2 \\
 \frac{\quad}{S/N : \lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x))}
 \end{array}$$

- eliminação de \:

$$\begin{array}{c}
 \underbrace{S/N}_{\downarrow} \quad : \quad \underbrace{\lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x))}_{\downarrow} \qquad \underbrace{(S/N) \quad \backslash \quad S}_{\downarrow} \quad : \quad \underbrace{\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))}_{\downarrow} \\
 \hline
 Y \quad : \quad \alpha \qquad + \qquad Y \quad \backslash \quad X \quad : \quad \phi \\
 \hline
 \Downarrow \\
 \hline
 X \quad : \quad (\qquad \phi \qquad \alpha \qquad) \\
 \downarrow \quad | \quad | \qquad \downarrow \qquad \downarrow \qquad | \\
 S \quad : \quad (\quad \underbrace{\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))}_{\downarrow} \quad \underbrace{\lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x))}_{\downarrow} \quad)
 \end{array}$$

- incluindo no diagrama a eliminação de \:

$\frac{\text{todo homem}}{\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))} \quad D$	$\frac{\text{ama}}{S : ((A\ w)\ z)} \quad D$	$\frac{\text{uma mulher}}{(S/N)\backslash S : \lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))} \quad D$
	$\frac{((A\ w)\ z)}{\lambda z.((A\ w)\ z)} \quad I\backslash^1$	
	$\frac{N\backslash S : \lambda z.((A\ w)\ z)}{S : (\lambda Q.\forall x.((H\ x) \rightarrow (Q\ x))\ \lambda z.((A\ w)\ z))} \quad E/$	
	$\begin{aligned} &=_{red.\beta} \forall x.((H\ x) \rightarrow (\lambda z.((A\ w)\ z)\ x)) \\ &=_{red.\beta} \forall x.((H\ x) \rightarrow ((A\ w)\ x)) \end{aligned}$	
	$\frac{}{S/N : \lambda w.\forall x.((H\ x) \rightarrow ((A\ w)\ x))} \quad I/^2$	
		$\frac{}{S : (\lambda Q_1.\exists y.((M\ y) \wedge (Q_1\ y))\ \lambda w.\forall x.((H\ x) \rightarrow ((A\ w)\ x)))} \quad E\backslash$

- redução- β :

$$\begin{array}{c}
 (\quad \lambda \quad Q_1 \quad . \quad \underbrace{\exists y.((M \ y) \wedge (Q_1 \ y))} \quad \underbrace{\lambda w.\forall x.((H \ x) \rightarrow ((A \ w) \ x))} \quad) \\
 \begin{array}{ccccccc}
 | & | & \downarrow & | & & \downarrow & & \downarrow & & | \\
 \hline
 (\quad \lambda \quad v \quad . & & \varphi & & & & \alpha & &)
 \end{array} \\
 \Downarrow \\
 \boxed{\varphi[v \mapsto \alpha]} \\
 \downarrow \\
 \overbrace{\exists y.((M \ y) \wedge (\lambda w.\forall x.((H \ x) \rightarrow ((A \ w) \ x)) \ y))}
 \end{array}$$

- incluindo no diagrama a redução- β :

todo homem	ama	uma mulher
$\frac{}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} D$	$\frac{}{S : ((A \ w) \ z)} D$	$\frac{}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} D$
	$\frac{}{N \setminus S : \lambda z. ((A \ w) \ z)} I \setminus^1$	
	$\frac{}{S : (\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) \ \lambda z. ((A \ w) \ z))} E/$	
	$\begin{aligned} &=_{red.\beta} \forall x. ((H \ x) \rightarrow (\lambda z. ((A \ w) \ z) \ x)) \\ &=_{red.\beta} \forall x. ((H \ x) \rightarrow ((A \ w) \ x)) \end{aligned}$	
	$\frac{}{S/N : \lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x))} I/^2$	
		$\frac{}{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x)))} E \setminus$
		$\begin{aligned} &=_{red.\beta} \exists y. ((M \ y) \wedge (\lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x)) \ y)) \end{aligned}$

- subfórmula a ser reduzida:

$$\exists y.((M \ y) \ \wedge \ \underbrace{(\lambda w.\forall x.((H \ x) \rightarrow ((A \ w) \ x)) \ y)}_{\text{aqui}}))$$

- redução- β :

$$\begin{array}{ccccccc} (& \lambda & w & . & \underbrace{\forall x.((H \ x) \rightarrow ((A \ w) \ x))}_{\downarrow} & y &) \\ | & | & \downarrow & | & & \downarrow & | \\ \boxed{(& \lambda & v & . & \varphi & \alpha &) \Rightarrow \varphi[v \mapsto \alpha]} \\ & & & & & & \downarrow \\ & & & & & & \underbrace{\forall x.((H \ x) \rightarrow ((A \ y) \ x))} \end{array}$$

- resultado:

$$\underbrace{\exists y.((M \ y) \ \wedge \ \underbrace{(\lambda w.\forall x.((H \ x) \rightarrow ((A \ w) \ x)) \ y}_{\forall x.((H \ x) \rightarrow ((A \ y) \ x))})}_{\exists y.((M \ y) \wedge \forall x.((H \ x) \rightarrow ((A \ y) \ x)))}$$

- incluindo no diagrama a segunda redução- β :

todo homem	ama	uma mulher
$\frac{}{S/(N \setminus S) : \lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x))} D$	$\frac{}{S : ((A \ w) \ z)} D$	$\frac{}{(S/N) \setminus S : \lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y))} D$
	$\frac{}{N \setminus S : \lambda z. ((A \ w) \ z)} I \setminus^1$	
$\frac{}{S : (\lambda Q. \forall x. ((H \ x) \rightarrow (Q \ x)) \ \lambda z. ((A \ w) \ z))} E/$		
$\quad =_{red.\beta} \forall x. ((H \ x) \rightarrow (\lambda z. ((A \ w) \ z) \ x))$		
$\quad =_{red.\beta} \forall x. ((H \ x) \rightarrow ((A \ w) \ x))$		
$\frac{}{S/N : \lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x))} I/2$		
$\frac{}{S : (\lambda Q_1. \exists y. ((M \ y) \wedge (Q_1 \ y)) \ \lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x)))} E \setminus$		
$\quad =_{red.\beta} \exists y. ((M \ y) \wedge (\lambda w. \forall x. ((H \ x) \rightarrow ((A \ w) \ x)) \ y))$		
$\quad =_{red.\beta} \exists y. ((M \ y) \wedge \forall x. ((H \ x) \rightarrow ((A \ y) \ x)))$		

4 Conclusão

- duas fórmulas diferentes:
 - $\forall x.((H\ x) \rightarrow \exists y.((M\ y) \wedge ((A\ y)\ x)))$
 - $\exists y.((M\ y) \wedge \forall x.((H\ x) \rightarrow ((A\ y)\ x)))$
- uma para cada interpretação
- ambas decorrentes unicamente dos procedimentos de análise

Referências

- [1] Bob Carpenter. *Type-Logical Semantics*. The MIT Press, Cambridge, MA, 1997.